# Gentzen's sequent calculus

1935.a. Gerhard Genzen defined 1st order formulas as sequents.

Sequent $A_1, \ldots, A_m \vdash B_1, \ldots, B_n$ is equivalent to 1st order formula $A_1 \land \ldots \land A_m \rightarrow B_1 \lor \ldots \lor B_n$ 

where  $m, n \ge 0$  and  $A_1, \dots, A_m, B_1, \dots, B_n$  are formulas

Sequent Ih formula  $A_1, ..., A_m$  – antecedent, rh formula  $B_1, ..., B_n$  – succedent.

Antecedent:  $A_1, ..., A_m$  represents formula  $A_1 \wedge ... \wedge A_m$ Succedent:  $B_1, ..., B_n$  represents formula  $B_1 \vee ... \vee B_n$ .

m = 0, means that rh formula is unconditionally true n = 0, means empty disjunct and contradiction

## Language

Sequents consist of formulas constructed using :  $\neg$ ,  $\land$ ,  $\lor$ ,  $\Rightarrow$ ,  $\forall$ ,  $\exists$ 

**Axiom** (scheme)  $A \rightarrow A$ Derivation rules - elimination and structural rules.

Each rule formalizes some proof step

### **Notations**

Upper case latin letters denote formuli

*x* – bound variable

a – free variable

t – term

 $\Gamma$ ,  $\Phi$ ,  $\Lambda$ ,  $\Pi$ – conjunctive/disjunctive sequences of formuli

## Inference rules I

#### **Axiom**

$$\frac{}{A \vdash A}$$
 (I)

$$\frac{}{A \vdash A} \quad (I) \qquad \qquad \frac{\Gamma \vdash \Delta, A \qquad A, \Sigma \vdash \Pi}{\Gamma, \Sigma \vdash \Delta, \Pi} \quad (Cut)$$

$$\frac{\Gamma, A \vdash \Delta}{\Gamma, A \land B \vdash \Delta} \quad (\land L_1)$$

$$\frac{\Gamma, A \vdash \Delta}{\Gamma, A \land B \vdash \Delta} \quad (\land L_1) \qquad \qquad \frac{\Gamma \vdash A, \Delta}{\Gamma \vdash A \lor B, \Delta} \quad (\lor R_1)$$

$$\frac{\Gamma, B \vdash \Delta}{\Gamma, A \land B \vdash \Delta} \quad (\land L_2)$$

$$\frac{\Gamma \vdash B, \Delta}{\Gamma \vdash A \lor B, \Delta} \quad (\lor R_2)$$

## Inference rules II

$$\frac{\Gamma, A \vdash \Delta \qquad \Sigma, B \vdash \Pi}{\Gamma, \Sigma, A \lor B \vdash \Delta, \Pi} \quad (\lor L) \qquad \frac{\Gamma \vdash A, \Delta \qquad \Sigma \vdash B, \Pi}{\Gamma, \Sigma \vdash A \land B, \Delta, \Pi} \quad (\land R)$$

$$\frac{\Gamma \vdash A, \Delta \qquad \Sigma \vdash B, \Pi}{\Gamma, \Sigma \vdash A \land B, \Delta, \Pi} \quad (\land R)$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma, \Sigma, A \to B \vdash \Delta, \Pi} \quad (\to L) \qquad \qquad \frac{\Gamma, A \vdash B, \Delta}{\Gamma \vdash A \to B, \Delta} \quad (\to R)$$

$$\frac{\Gamma, A \vdash B, \Delta}{\Gamma \vdash A \to B, \Delta} \quad (\to R)$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma, \neg A \vdash \Delta} \quad (\neg L) \qquad \qquad \frac{\Gamma, A \vdash \Delta}{\Gamma \vdash \neg A, \Delta} \quad (\neg R)$$

$$\frac{\Gamma, A \vdash \Delta}{\Gamma \vdash \neg A, \Delta} \quad (\neg R)$$

## Inference rules III

$$\frac{\Gamma, A[t/x] \vdash \Delta}{\Gamma, \forall x A \vdash \Delta} \quad (\forall L)$$

$$\frac{\Gamma \vdash A[y/x], \Delta}{\Gamma \vdash \forall x A, \Delta} \quad (\forall R)$$

$$\frac{\Gamma, A[y/x] \vdash \Delta}{\Gamma, \exists x A \vdash \Delta} \quad (\exists L)$$

$$\frac{\Gamma \vdash A[t/x], \Delta}{\Gamma \vdash \exists x A, \Delta} \quad (\exists R)$$

In rules  $\forall R$  and  $\exists L$  the variable y must not occur free within  $\Gamma$  and  $\Delta$ . Alternatively, the variable y must not appear anywhere in the respective lower sequents.

### Structural rules

$$\frac{\Gamma \vdash \Delta}{\Gamma, A \vdash \Delta} \quad (WL)$$

$$\frac{\Gamma \vdash \Delta}{\Gamma \vdash A, \Delta} \quad (WR)$$

$$\frac{\Gamma, A, A \vdash \Delta}{\Gamma, A \vdash \Delta} \quad (CL)$$

$$\frac{\Gamma \vdash A, A, \Delta}{\Gamma \vdash A, \Delta} \quad (CR)$$

$$\frac{\Gamma_1, A, B, \Gamma_2 \vdash \Delta}{\Gamma_1, B, A, \Gamma_2 \vdash \Delta} \quad (PL)$$

$$\frac{\Gamma \vdash \Delta_1, A, B, \Delta_2}{\Gamma \vdash \Delta_1, B, A, \Delta_2} \quad (PR)$$

# Inference Example

$$\frac{B \vdash B}{B \lor C \vdash B, C} (I) \\ \hline B \lor C \vdash B, C \\ \hline B \lor C \vdash C, B \\ \hline B \lor C \vdash C, B \\ \hline (\neg L) \\ \hline (B \lor C), \neg C, (B \to \neg A) \vdash \neg A \\ \hline (B \lor C), \neg C, ((B \to \neg A) \vdash \neg A) \\ \hline (B \lor C), \neg C, ((B \to \neg A) \land \neg C) \vdash \neg A \\ \hline (B \lor C), ((B \to \neg A) \land \neg C), \neg C \vdash \neg A \\ \hline (B \lor C), ((B \to \neg A) \land \neg C), \neg C \vdash \neg A \\ \hline (B \lor C), ((B \to \neg A) \land \neg C), ((B \to \neg A) \land \neg C) \vdash \neg A \\ \hline (B \lor C), ((B \to \neg A) \land \neg C), ((B \to \neg A) \land \neg C) \vdash \neg A \\ \hline (B \lor C), ((B \to \neg A) \land \neg C), (B \lor C) \vdash \neg A \\ \hline (B \lor C), ((B \to \neg A) \land \neg C), (B \lor C) \vdash \neg A \\ \hline (B \to \neg A) \land \neg C), (A \to (B \lor C)) \vdash \neg A, \neg A \\ \hline (B \to \neg A) \land \neg C), (A \to (B \lor C)) \vdash \neg A \\ \hline (A \to (B \lor C)), ((B \to \neg A) \land \neg C) \to \neg A \\ \hline (A \to (B \lor C)) \vdash (((B \to \neg A) \land \neg C) \to \neg A) \\ \hline (A \to (B \lor C)) \vdash (((B \to \neg A) \land \neg C) \to \neg A) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A)) \\ \hline (A \to (B \lor C)) \to (((B \to \neg A) \land \neg C) \to \neg A))$$