ITC8190 Mathematics for Computer Science Binary Relations on a Set

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A binary relation R on a set A is the subset

 $R \subseteq A \times A : xRy \iff (x, y) \in R$.

The relation < on a set $A = \{1, 2, 3\}$ is the subset $\{(1, 2), (1, 3), (2, 3)\}.$

Relation R on a set A is **reflexive** if every element x in A is related to itself. It means that

$$\forall x \in A : xRx$$
.

Example: the relation \leq on \mathbb{Z} is reflexive, but the relation < is not.

R is called **anti-reflexive** if every element x in A is not related to itself.

$$\forall x \in A : \neg(xRx)$$
.

Relation < on \mathbb{Z} is anti–reflexive.

Relation R on a set A is called **symmetric** if for any pair of elements x, y in A, it holds that if x is related to y, then y is related to x.

 $\forall x, y \in A : xRy \implies yRx$.

Example: the relation = on \mathbb{R} is symmetric, since for all $a, b \in \mathbb{R}$ it holds that a = b implies b = a.

Relation R on a set A is **anti-symmetric** if for any pair of elements x, y in A it holds that if x is related to y, and y is related to x, then x and y are the same element (written as x = y).

$$\forall x, y \in A : xRy \land yRx \implies x = y .$$

Example: relation \leq is anti-symmetric, since

$$x \leqslant y \land y \leqslant x \implies x = y$$
.

Relation R on a set A is **asymmetric** if it holds that if x is related to y, then y is unrelated to x.

$$\forall x, y \in A : xRy \implies \neg(yRx)$$
.

Example: the relation < on \mathbb{R} is asymmetric, and the condition x < y implies that $y \not< x$.

$$x < y \implies \neg (y < x)$$
.

Relation R on a set A is **transitive** if

$$\forall x, y, z \in A : xRy \land yRz \implies xRz$$
.

Example: relations < and = are transitive. It can be seen that

$$a < b \land b < c \implies a < c$$
,
 $a = b \land b = c \implies a = c$.

Symmetric and transitive relation is reflexive.

Proof.

By symmetry,

$$xRy \implies yRx$$
.

By transitivity,

$$xRy \wedge yRx \implies xRx$$
.

Therefore, symmetry and transitivity imply reflexivity.

Asymmetric relation is anti-reflexive.

Proof.

By asymmetry, $xRy \Longrightarrow \neg(yRx)$. Since y can be any element, let y = x. Then $xRx \Longrightarrow \neg(xRx)$. Hence, asymmetry implies anti-reflexivity.

Anti-reflexive and transitive relation is asymmetric.

Proof.

Indeed, it can be seen that $xRy \wedge yRx$ is always false. By transitivity,

$$xRy \wedge yRx \implies xRx$$
,

which contradicts with anti–reflexivity. So xRy and yRx cannot happen at the same time. Therefore,

$$xRy \implies \neg(yRx)$$
.

Anti-reflexive and transitive relation is anti-symmetric.

Proof.

By transitivity,

$$xRy \wedge yRx \implies xRx$$
,

which contradicts with the anti–reflexivity property. And so, the implication

$$xRy \wedge yRx \implies x = y$$

is true.

Corollary 1

If the relation is anti-reflexive and transitive, then anti-symmetry is the same as asymmetry.

Anti-reflexive relation is anti-symmetric iff it is asymmetric.

Proof.

First, we show that if anti-reflexive relation is asymmetric, then it is anti-symmetric. We need to show that $xRy \wedge yRx \implies x=y$. By assymetry, $xRy \implies \neg(yRx)$, the expression $xRy \wedge yRx$ is always false, and hence the implication $xRy \wedge yRx \implies x=y$ is true.

Secondly, we show that if anti-reflexive relation is anti-symmetric, then it is asymmetric. We need to show that $xRy \implies \neg(yRx)$. Let xRy. If yRx is true, then by anti-symmetry $xRy \land yRx$ would imply x = y. If yRx is true and y = x, then xRx is true. A contradiction with anti-reflexivity. And so, if xRy is true, yRx must be false. Hence $xRy \implies \neg(yRx)$.

Relation R on a set A is **connex** if any pair of elements in A is comparable under R.

$$\forall x, y \in A : xRy \vee yRx$$
.

R is called **trichotomous** if any pair of elements in A is either comparable under R or is the same element.

$$\forall x, y \in A : xRy \veebar yRx \veebar x = y .$$

