

Exercise 1

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{N>0 ∧ N=n}                                % ≡ Pre
Z:=1
{Pre ∧ Z=1}                                  % annotation
WHILE N > 0 DO
    {N ≥ 0 ∧ Z * M^N = M^n}                % ≡ Inv
    BEGIN
        Z := Z * M;
        N := N - 1;
    END;
{Z= M^n}                                     % ≡ Post

```

Exercise 2

Prove partial correctness of the guarded command program specification:

$$\{x \geq 0\} \quad y := 1; \star[y \times y \leq x \rightarrow y := y + 1]; y := y - 1 \quad \{y^2 \leq x < (y + 1)^2\}$$

Hint: try with invariant $(y - 1)^2 \leq x$

Exercise 3

Given an annotated program $S_1||S_2$ verify if it is interference free and prove the partial correctness of S_1 and S_2 separately.

$$\begin{aligned}
P_1 &\equiv \{x \leq 4 \wedge y = 2\} \\
&\quad S_1: \langle x \geq 2 \rightarrow y := y - 2 \rangle \\
Q_1 &\equiv \{y \leq x \wedge x \geq 0\} \\
&\quad \| \\
P_2 &\equiv \{x \geq 0 \wedge y \geq 0\} \\
&\quad S_2: \langle x = 4 \wedge y = 1 \rightarrow z := x - 3 \rangle \\
Q_2 &\equiv \{y + 2 \leq x\}
\end{aligned}$$

Exercise 4

Specify the cooperation tests for channels C and D

$$\begin{aligned}
P &\equiv \{x = 6 \wedge u = 0 \wedge y - x = 6\} \\
P_1 &\equiv \{x \geq 5 \wedge y > 7\} \\
&\quad S_1: \langle C! x + 3 \rangle; \{x > 0 \wedge y \geq 7\} \langle D? y \rangle; \{y > 10 \wedge x > 0\} \\
Q_1 &\equiv \{y > 9 \wedge x > 0\} \\
&\quad \| \\
P_2 &\equiv \{u = 0\} \\
&\quad S_2: \langle C? u \rangle; \{u = 8\} \langle D! u+5 \rangle \\
&\quad \square \\
&\quad S_2: \langle C? u \rangle; \langle u := u - 1 \rangle \{u = 7\} \langle D! u+5 \rangle \\
Q_2 &\equiv \{u < 9\} \\
Q &\equiv \{x > 0 \wedge u < 10 \wedge y > 0\}
\end{aligned}$$

Exercise 5

Let R and T be nonempty sets of natural numbers. Consider the following partitioning algorithm $S_1 \parallel S_2$, where

$S_1 \equiv max := max(R); c?mn; d!max;$
 $\star [max > mn \rightarrow R := (R \setminus \{max\}) \cup \{mn\}; max := max(R);$
 $c?mn; d!max]$

$S_2 \equiv min := min(T); c!min; d?mx;$
 $\star [mx > min \rightarrow T := (T \setminus \{min\}) \cup \{mx\}; min := min(T);$
 $c!min; d?mx]$

Prove, by means of the method of Levin & Gries,

$\{R = R_0 \neq \emptyset \wedge T = T_0 \neq \emptyset \wedge R \cap T = \emptyset\} \quad S_1 \parallel S_2$

$\{|R| = |R_0| \wedge |T| = |T_0| \wedge R \cup T = R_0 \cup T_0 \wedge max(R) < min(T)\}$
where, for a set A , $|A|$ denotes the number of elements of A , and R_0 and T_0 are logical variables denoting a finite set of natural numbers.