Lecture 3
Property specification in Temporal Logic CTL*

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Model Checking

$$M \models P\ ?$$

Given

- M – model
- P – property to be checked
- $\models$ – satisfiability relation („*M satisfies P*“)

Check if *M satisfies P*

Model: Kripke Structure (revisited I)

- KS is a state-transition system that captures
  - what is true in a state (denoted as labeling of the state)
  - what can be viewed as an atomic move (denoted as transition)
  - the succession of states   (paths on the model graph)

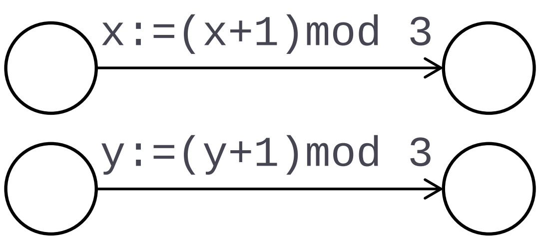
- KS is a static representation that can be unfolded to a *tree of execution traces* on which temporal properties are verified.

Representing transition as formuli

- In Kripke structure, transition $(s, s') \in R$ corresponds to one step of executing the program.

- Suppose a program has two steps
  - `x := (x+1) mod 3;`
  - `y := (y+1) mod 3.`

  Then $R = \{R_1, R_2\}$
  - $R_1 : (x' = (x+1)\ mod\ 3) \wedge (y' = y)$
  - $R_2 : (y' = (y+1)\ mod\ 3) \wedge (x' = x)$



Consecutive States

- State space:

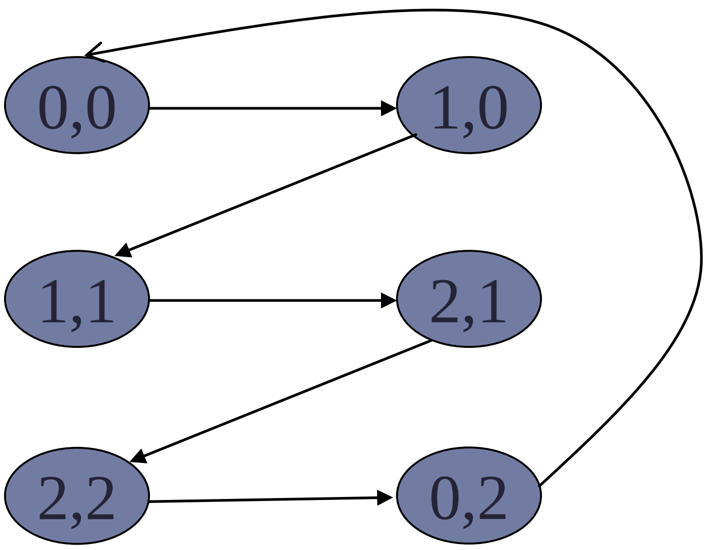
we can restrict our attention to pairs of consecutive states $s = (x, y)$ and $s' = (x', y')$ on the state space $\{0, 1, 2\} \times \{0, 1, 2\}$, i.e.

$$s, s' \in \{0, 1, 2\} \times \{0, 1, 2\}$$

- Question: Can we construct a logic formula that describes the relation between <u>any</u> two consecutive states s and s'?

- Individual consecutive states can be related by R_1 or R_2.

Consecutive states represented by $R_1 \vee R_2$



Representing transitions (revisited II)

- In Kripke structure, a transition (s, s') $\in R$ corresponds to one step of execution of the program

- Suppose a program P has two steps
  - `x := (x+1) mod 3;`
  - `y := (y+1) mod 3;`

- For the whole program we have
  $R = ((x' = x+1 \textit{ mod } 3) \wedge y' = y) \vee ((y' = y+1 \textit{ mod } 3) \wedge x'=x)$

- (s, s') that satisfies R means "from state s we can get to s' by any step of execution that satisfies R"

A giant *R*

- We can compute *R* for the whole program
  - then we will know whether any two states are one-step reachable
- Convenient, but globally we loose information:

  e.g., the order in which the statements are executed
- Comment:
  - without order, the disjuncts have <u>no clear precedence</u>!

Introducing program counter

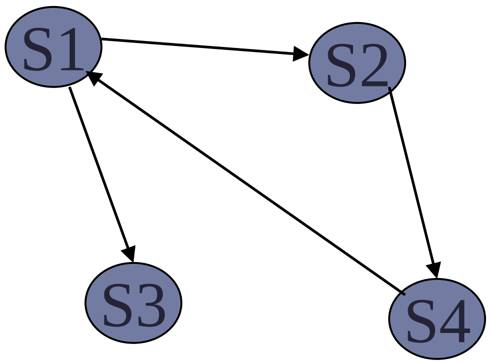
- In a real machine, the order of execution is managed by *program counters.*
- We introduce a virtual variable pc, and assume the program commands are labeled with $l_0 \ldots l_n$.
- For instance
  - In the program:
    - `l0: x := x+1;`
    - `l1: y := x+1;`
    - `l2: …`
  - In the logic:
    - $R_1 : x' = x+1 \land y'=y \land pc = l_0 \land pc' = l_1$
    - $R_2 : y' = y+1 \land x'=x \land pc = l_1 \land pc' = l_2$

Now we have complete logic representation of program executions in our model *M*!

Temporal logic CTL*

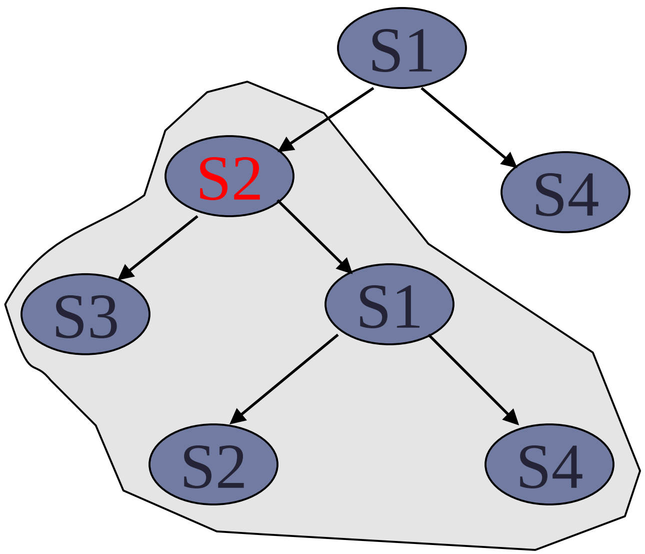
- Semantics

  KS and its logic representation are static models of program execution



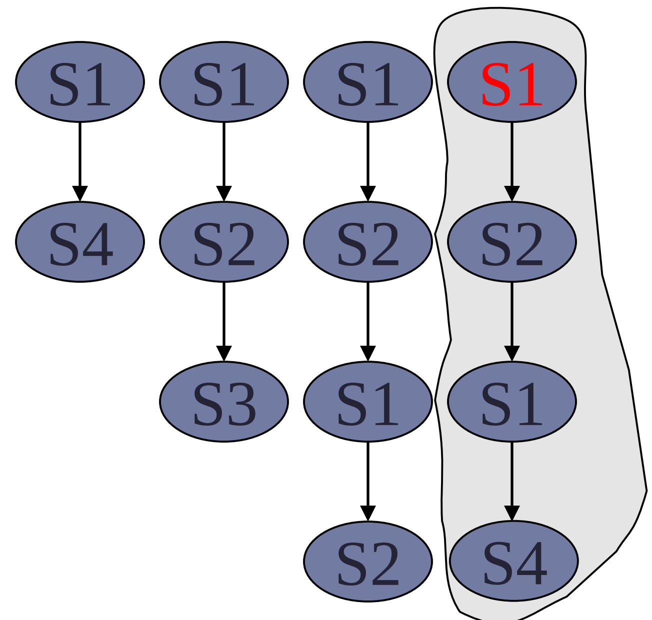
Dynamic model of program execution = unfolding of the static model

Tree structure: branching time



Is a formula valid at a given node, which represents a subtree?

Traces: linear time



Is a formula valid along a given path?

CTL* (Computational Tree Logic)

- Combines branching time and linear time
- Basic Operators
  - X: neXt
  - F: Future (〈〉)
  - G: Global ([])
  - U: Until
  - R: Release

CTL*

- State formulas
  - Express a propery of a state
  - Path quantifiers:
    - **A** – for all paths, **E** – for some paths
- Path formulas
  - Expess a property of a path
  - State quantifiers:
    - **G** – for all states (of the path)
    - **F** – for some state (of the path)

State Formulas (1)

- Atomic propositions
  - $p \in \mathrm{AP}$, then p is a state formula
  - Examples: $x > 0$, $odd(y)$
- Propositional combinations of state formulas
  - $\neg\, \varphi, \quad \varphi \vee \psi, \quad \varphi \wedge \psi \ldots$
  - Examples:
    - $x > 0 \vee odd(y)$,
    - $req \Rightarrow (\mathrm{AF}\ ack)$
      - "A" is a path quantifier
      - "F ack" is a path formula
      - "AF ack" is a state formula (interpreted in a state)

State Formulas (2)

- Quantifiers A and E construct a state formula from a path formula

- $E\varphi$, where φ is a path formula, which expresses property of a path
  - E means “there exists”
  - E φ - on some path from this state on φ is *true*.

- A φ
  - A means “for all paths“
  - A φ - on all paths starting from this state φ is *true*.

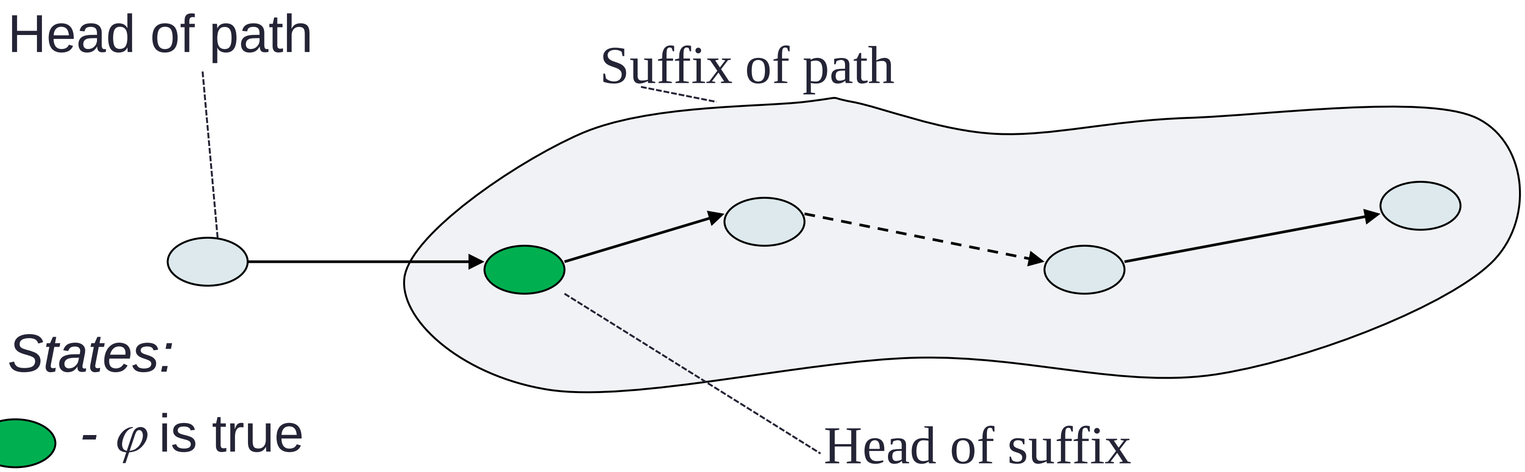
Forms of Path Formulas

- A state formula φ
  - φ is true for the <u>first state</u> of this path
- For path formulas φ and ψ, the path formulas are:
  - $\neg \varphi$, $\varphi \vee \psi$, $\varphi \wedge \psi$
  - X φ, Fφ, G φ, φ U ψ, φ R ψ
    - *X – next*
    - *F – eventually*
    - *G – globally*
    - *U – until*
    - *R – release*

Path Formulas (I): *Next*-operator

X φ, where φ is a path formula

- φ is valid for the suffix of this path (path minus the first state)



States:

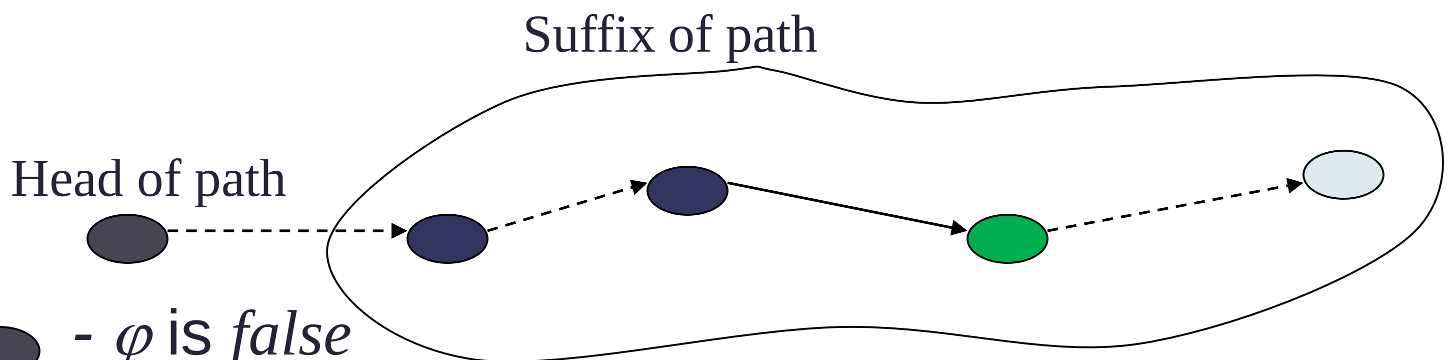
- φ is true

- φ can be either true or false in other states

Path Formulas II: *Eventually*-operator

F φ:

φ is valid for this path



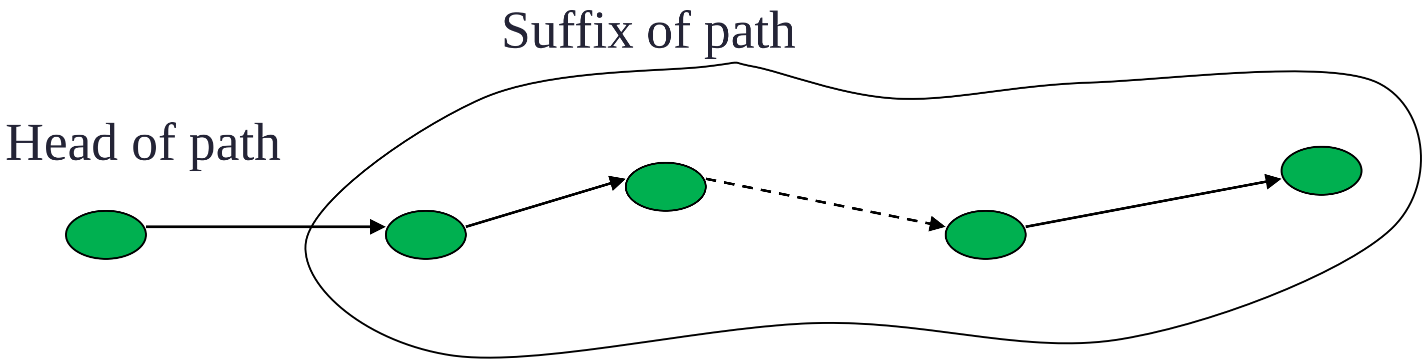
- φ is *false*

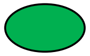
- φ is *true*

- φ can be either *true* or *false*

Path Formulas (III): *Globally*-operator

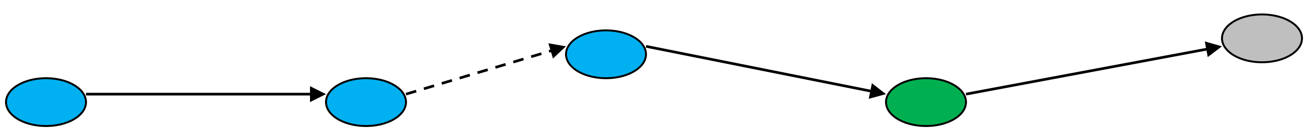
- G φ
  - φ is valid for head and every suffix of this path

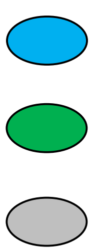



- φ is true

Path Formulas IV: *Until*-operator

- $\varphi \, \mathrm{U} \, \psi$
  - ψ is valid on a suffix of the path, before the first node of which φ is valid on every suffix thereon





- φ is true

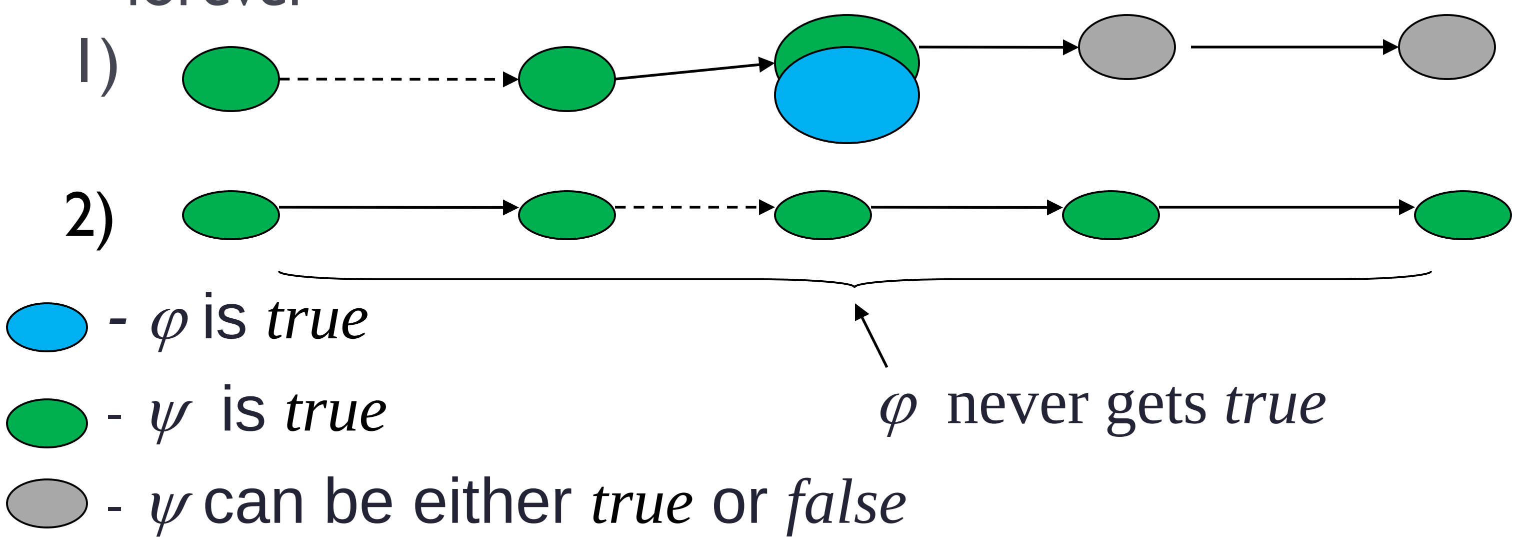
- ψ is true

- φ and ψ are either true or false

Path Formulas (V): *Release*-operator

φ R ψ

- ψ has to be *true* until and including the point where φ becomes *true*; if never becomes *true* must remain *true* forever



Formal semantics of CTL* (1)

- Notations
  - $M, s \vDash \varphi$ iff φ holds in state s of model M
  - $M, \pi \vDash \varphi$ iff φ holds along the path π in M
  - π^i : i-th suffix of π
    - $\pi = s_0, s_1, \ldots$, then $\pi^1 = s_1, \ldots$

Semantics of CTL* (2)

- *Path formulas are interpreted over a path*:
  - $M, \pi \models \varphi$
  - $M, \pi \models \mathrm{X}\ \varphi$
  - $M, \pi \models \mathrm{F}\ \varphi$
  - $M, \pi \models \varphi\, \mathrm{U}\, \psi$

Semantics of CTL* (3)

- *State formulas are interpreted over a set of states (of a path)*
  - $M, s \models p$
  - $M, s \models \neg \varphi$
  - $M, s \models \mathrm{E}\ \varphi$
  - $M, s \models \mathrm{A}\ \varphi$

CTL vs. CTL*

- CTL*, CTL and LTL have different expressive powers:
- Example:
  - In CTL there is no formula being equivalent to LTL formula A(FG p).
  - In LTL there is no formula equivalent to CTL formula AG(EF p).
  - $A(FG\ p) \vee AG(EF\ p)$ is a CTL* formula that cannot be expressed neither in CTL nor in LTL.

CTL

- Quantifiers over paths
  - A – **A**ll: has to hold on all paths starting from the current state.
  - E – **E**xists: there exists at least one path starting from the current state where holds.

- In CTL, path formulas can occur only when paired with A or E , *i.e.* one path operator followed by a state operator.

  if φ and ψ are path formulas, then
  - X φ,
  - F φ,
  - G φ,
  - φ U ψ,
  - φ R ψ

  are path formulas

LTL (contains only path formulas)

Form of path formulas:

- If $p \in \mathrm{AP}$, then p is a path formula
- If φ and ψ are path formulas, then
  - $\neg\varphi$
  - $\varphi \vee \psi$
  - $\varphi \wedge \psi$
  - $\mathrm{X}\ \varphi$
  - $\mathrm{F}\ \varphi$
  - $\mathrm{G}\ \varphi$
  - $\varphi\, \mathrm{U}\, \psi$
  - $\varphi\, \mathrm{R}\, \psi$

are path formulas.

Minimal set of CTL temporal operators

- Transformations used for temporal operators :
  - $EF\ \varphi == E\ [true\ U\ \ \varphi\ ]$ (because $F\ \varphi == [true\ U\ \varphi\ ]$)
  - $AX\ \varphi == \neg\ EX(\neg\ \varphi\ )$
  - $AG\ \varphi == \neg\ EF(\neg\ \varphi\ ) == \neg\ E\ [true\ U\ \varphi\ ]$
  - $AF\ \varphi == A\ [true\ U\ \varphi\ ] == \neg\ EG\ \neg\ \varphi$
  - $A[\varphi\ U \psi] == \neg(\ E[(\neg\ \psi)\ U\ \neg(\varphi \vee\ \psi)] \vee EG\ (\neg \psi)\ )$

Summary

- CTL* is a general temporal logic that offers strong expressive power, more than CTL and LTL separately.
- CTL and LTL are practically useful enough; CTL* helps us to understand the relations between LTL and CTL.
- Next we will show how to check satisfiability of CTL formuli on Kripke structures