## 1 Exercises

**Exercise 1.** How would you test if an integer n is prime, using the brute-force approach?

**Solution.** We need to check if n is co-prime to all integers in the range  $[2, \sqrt{n}]$ . If it is the case, then n is prime.

Exercise 2. Apply Fermat primality test to verify the primality of 351, 511, and 717.

**Solution.** For 351, take a = 2 and we have  $2^{350} \mod 351 = 121 \neq 1$ , hence, 2 is a Fermat witness, and 351 is composite.

For 511, take a = 2 and we have  $2^{510} \mod 511 = 64 \neq 1$ , hence 2 is a Fermat witness, and 511 is composite.

For 717, take a = 2 and we have  $2^{716} \mod 717 = 4 \neq 1$ , hence 2 is a Fermat witness, and 717 is composite.

**Exercise 3.** Solve for x:  $x^2 \mod 63 = 1$ , where 63 is the product of two primes 7 and 9.

Solution. We need to find the square roots of 1 is  $\mathbb{Z}_{63}$ . By the Chinese Remainder Theorem,  $\mathbb{Z}/63\mathbb{Z} \cong \mathbb{Z}/7\mathbb{Z} \times \mathbb{Z}/9\mathbb{Z}$ , and isomorphism  $\varphi : \mathbb{Z}/63\mathbb{Z} \to \mathbb{Z}/7\mathbb{Z} \times \mathbb{Z}/9\mathbb{Z}$  is given by  $\varphi(n \mod 63) \mapsto (x \mod 7, x \mod 9) = (x_7, x_9)$ . To find the square roots of unity, we need to find elements  $(x_7, x_9)$  that would satisfy  $(x_7^2 \mod 7, x_9^2 \mod 9) = (1, 1)$ . For this, we solve a system of equations  $\begin{cases} x_7^2 \mod 7 = 1 \ x_9 \mod 9 = 1 \ x_9 \mod 7 = 1 \ x_9 \mod 7 = 1 \ x_9 \mod 9 = 8 \ x_9^2 \mod 9 = 1 \ x_9 = 8 \ x_{10} \mod 10 \ x_{10} = 10 \ x_{10} \ x_{10} (x_{10}, x_{10}) \ x_{10} (x_{10}, x_{10}) \ x_{10} (x_{10}, x_{10}) \ x_{10} (x_{10}, x_{10}) \ x_{10} \oplus 10 \ x_{10} \$ 

To conclude, there are four 2-nd roots of unity in  $\mathbb{Z}/63\mathbb{Z}$ , they are: 1, 8, 55, 62.

**Exercise 4.** Adversary Carol has intercepted two RSA cryptograms,  $y_1 = 537$  sent by Alice to Bob, and  $y_2 = 285$  sent by Alice to Eve. The public key of Bob is  $(e_1 = 18, n_1 = 943)$ , and the public key of Eve is  $(e_2 = 19, n_2 = 943)$ . What is the message *m* sent by Alice to Bob and Eve?

Solution. Carol knows that Bob and Eve receive the same message m. Hence,

$$537 = m^{18} \mod 943 \ , \\ 285 = m^{19} \mod 943 \ .$$

Since the public exponents of Bob and Eve are co-prime, meaning that gcd(18, 19) = 1, the Bézout identity applies.

$$gcd(18, 19) = 1 \cdot 19 + (-1) \cdot 18 = 1$$
.

Carol needs to combine  $y_1$  and  $y_2$  in such a way that would exploit the Bézout identity to get m. Observe that given  $\alpha \cdot e_1 + \beta \cdot e_2 = 1$  implies that

$$y_1^{\alpha} \cdot y_2^{\beta} = (m^{e_1})^{\alpha} \cdot (m^{e_2})^{\beta} = m^{\alpha \cdot e_1} \cdot m^{\beta \cdot e_2} = m^{\alpha \cdot e_1 + \beta \cdot e_2} = m^1 = m$$
.

Given that  $\alpha = -1$ , and hence  $y_1^{\alpha} = 537^{-1} \equiv 72 \pmod{943}$  and  $\beta = 1$ , Carol needs to compute

 $m = 72 \cdot 285 \mod 943 = 717$ .

**Exercise 5.** Suppose that adversary Carol has intercepted three cryptograms  $y_1, y_2, y_3$  sent to three different users whose public keys are  $(e = 3, n_1), (e = 3, n_2), (e = 3, n_3)$ . Notice that all public keys use the same public exponent, and different moduli. What does Carol needs to do to reconstruct the message m?

Solution. Carol knows that

$$y_1 = m^3 \mod n_1$$
$$y_2 = m^3 \mod n_2$$
$$y_3 = m^3 \mod n_3$$

and she knows that  $m < n_1, m < n_2, m < n_3$ . Hence,  $m^3 < n_1 \cdot n_2 \cdot n_3$ , and therefore  $m^3$  is the solution to the CRT

$$\begin{cases} x \mod n_1 = y_1 \\ x \mod n_2 = y_2 \\ x \mod n_3 = y_3 \end{cases}$$

CRT guarantees the existence of the solution, and that the solution is unique. Since  $m^3$  solves the CRT, it must be this unique solution. All that Carol needs to do now is to calculate

$$m = \sqrt[3]{m^3}$$

to reconstruct the message m.

If the public exponent is relatively small, say, n, then the adversary needs to intercept n cryptograms in order to be able to decipher it.